Monday, June 17, 2024

Exercise 1

Remember that by "vertex coloring" of an undirected graph we mean the task of assigning an element from a finite set of labels ("colors") to each vertex of the graph so that no two connected vertices have the same color; we say that a graph can be k-colored if the task can be successfully carried out with no more than k colors.

For each of the following coloring-related languages, determine whether they are in **P**, **NP** and/or **coNP** and provide a short motivation for your answers:

- 1. graphs that can be 2-colored;
- 2. graphs that cannot be 2-colored;
- 3. graphs that can be 3-colored;
- 4. graphs that cannot be 3-colored;
- 5. graphs G = (V, E) that cannot be k-colored for any k < |V|.

Acceptable motivations can be formulated as short descriptions of algorithms (verbal or any kind of pseudocode), set inclusions (e.g., "I already proved that $L \in A$, but we know that $A \subset B$, therefore $L \in B$ "), reductions to some **NP**-complete version of SATISFIABILITY (e.g., 3-SAT).

Exercise 2

Choose one of the following topics and develop it along the suggested lines.

Topic 1: Kolmogorov complexity

- 1. Given a suitable alphabet Σ and a universal Turing machine (UTM) \mathcal{U} , define the Kolmogorov complexity $K_{\mathcal{U}}(s)$ of a string $s \in \Sigma^*$.
- 2. Show that, given two UTMs \mathcal{U} and \mathcal{V} , the difference between the Kolmogorov complexities $K_{\mathcal{U}}(s)$ and $K_{\mathcal{V}}(s)$ of the same string s is bounded by a constant that only depends on \mathcal{U} and \mathcal{V} and not on s.
- 3. Outline the proof that $K_{\mathcal{U}}$ is uncomputable.

Topic 2: Post Correspondence Problem

- 1. Define the Post Correspondence Problem (PCP).
- 2. Given the "Restricted" version of the Post Correspondence Problem (RPCP) where the initial pair is fixed (i.e., pre-set as part of the instance), show how an instance of RPCP can be reduced to an equivalent instance of PCP.
- 3. Outline the proof that PCP is uncomputable.